# 4. Counting faces

#### MAIN QUESTIONS:

- thow many  $\delta$ -dimensional facer can a polytope have for some  $\delta \in \{0,...,d-1\}$ ?
- · What relations exist between these face numbers?
- face numbers of polytopes?

Def: The list of numbers usually not included 
$$\overrightarrow{f} = \overrightarrow{f}(P) = (f_{-1}, f_0, f_1, \dots, f_{d-1}, f_d)$$

$$= 1$$

with  $f_{\delta} = f_{\delta}(P) := \# faces of dimension <math>\delta$  is called the f-vector of P.

Eg.: 2D:  $f_0 = f_1$  and  $f_0, f_1 \ge 3$ completely characterizes f-vectors in dimension d = 2.

# 4.1. Euler's polyhedral formula

# and some consequences

$$V - E + F = 2 \qquad \leftarrow \text{ holds for all}$$
or  $f_0 - f_1 + f_2 = 2$  connected planar
graphs

Ex: prove it by inductively adding vertices and edges.

• f-vectors of 3-polytopes have been completely characterized

Thm: (Steinitz)

 $f = (f_0, f_1, f_2)$  is an f-vector of a 3-polytope iff

(i) 
$$f_0 - f_1 + f_2 = 2$$
 (Euler)

(ii) 
$$4 \le f_0 \le 2f_2 - 4$$
 (iii)  $4 \le f_2 \le 2f_0 - 4$  dual

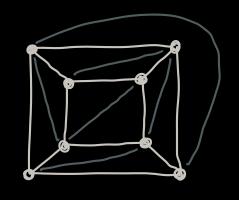
(iii) 
$$4 \leq f_2 \leq 2f_0 - 4$$

$$\begin{pmatrix} f_0 \\ f_1 \\ f_2 \end{pmatrix} \xrightarrow{P} \stackrel{P^{\circ}}{P} \begin{pmatrix} f_1 \\ f_0 \\ f_0 \end{pmatrix}$$

- · no proof, but I tell you where (ii) and (iii) come from:
  - try to maximize the number of edges and 2-faces by adding edges until you can't without making the graph non-planor

# -> graph becomes a triangulation

(= every 2-face is a triangle)



- let's count incident edge-face pairs in two ways

$$2f_1 = \sum_{e \text{ fige}} \sum_{f \text{ eff}} 1 = \sum_{g \text{ fige}} \sum_{g \text{ fige}} 1 = 3f_2$$

$$\longrightarrow f_1 = \frac{3}{2} f_2$$

(iii) 
$$\longrightarrow$$
  $f_2 \stackrel{\leq}{=} 2f_0 - 4$ 

(ii) follows from duality

since we maximized for this is on inequality for general planor graphs.

## Some consequences of Euler's polyhedral formula

• let's count incident vertex-edge pairs in general planar graphs:

$$2f_1 = \sum_{v} \delta(v) - \delta f_0 \qquad \longrightarrow \qquad f_0 = \frac{2f_1}{\delta}$$
Vortex degree average degree

· and incident face-edge pairs:

$$2f_1 = \sum_{f} g(f) = \overline{g} f_2 \qquad \rightarrow \qquad f_2 = \frac{2f_1}{\overline{g}}$$
gonality (number average gonality of vertices of a 2-face)

· plug into Euler's polyhedral formula

$$2 = \int_0 -f_1 + f_2 = \frac{2}{\delta} f_1 - f_1 + \frac{2}{g} f_1$$

$$\rightarrow \frac{1}{\delta} + \frac{1}{g} = \frac{1}{2} + \frac{1}{f_1} > \frac{1}{2}$$

• we know that  $\overline{\delta} \ge 3$  (Balinski)

- Cor: Every 3-polytope has a 2-fooe with at most five vertices
  - Every 3-polytope has a vertex of degree at most five.

    Ex: if all vertex-degrees

~ "Ramsey theory for polytopes" a triangular 2-face

Cor: Every polytope has a 2-face that is a triangle, quadrangle or pentason.

Proof: every polytope has a 3-face to which the previous corollary applies.

BUT: more is known!

Thm: (Kalai) Every d-polytope with  $d \ge 5$  has a 2-face that is either a triangle or quadranole.

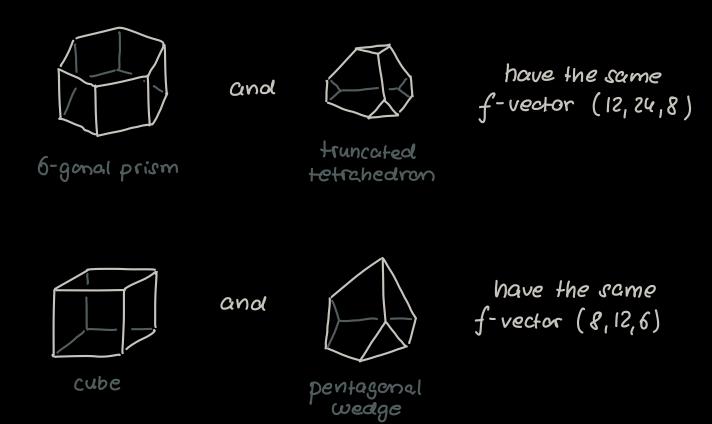
Note: the 4-dimensional 120-cell has only pentagonal 2-faces.

Can similar things be said about higher-dimensional foces?

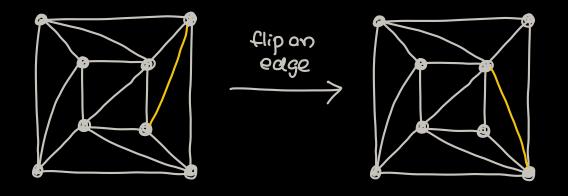
OPEN: Do all sufficiently high-climensional polytopes have a 3-face that is either a tetrahedron or a cube?

Interlude: is a polytope determined by its f-vector?

· No, already not true in dimension 3.



· Simple procedure for creating counterexamples:



### A glimpse at the Euler-Poinceré identity

Q: Does Euler's polyhedral formula generalize to higher climensions? And in which form?

The pattern becomes clearer if we include for end fol

3D 
$$-f_{-1} + f_0 - f_1 + f_2 - f_3 = 0$$
  
 $\rightarrow f_0 - f_1 + f_2 = f_{-1} + f_3 = 2$   
4D  $-f_{-1} + f_0 - f_1 + f_2 - f_3 + f_4 = 0$   
 $\rightarrow f_0 - f_1 + f_2 - f_3 = f_{-1} - f_4 = 0$ 

The right generalization seems to be

Thm: (Euler-Poincaré identity) For a d-polytope holds

$$\sum_{i=-1}^{d} (-1)^{i} f_{i} = 0 \quad \text{or equivalently} \quad \sum_{i=0}^{d-1} (-1)^{i} f_{i} = 1 - (-1)^{d}$$

( proof will be given next week; we need: shellability)

• This is the only linear relation that holds between the face-numbers of a general paytope!

## 4.2. The Dehn-Sommonville equations

• Much more can be said about the f-vector if the polytope is simple or simplicial

Thm: (Dehn-Sommerville equations)

PCRd simple (analogously for simplicial)

For all KE {0,..., d} holas

$$\sum_{i=k}^{d} (-1)^{i} {i \choose k} f_{i} = \sum_{i=d-k}^{d} (-1)^{d-i} {i \choose d-k} f_{i}$$

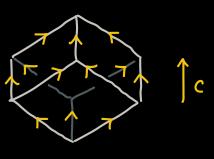
#### Remarks:

- might look scary but become much nicer after a basis chanse (see proof)
- · Dseq for k and al-k are equivalent
- · these are Ld/21 linearly independent equations
  - -> knowing only half the f-vector

    we can reconstruct the missing entries

#### Proof:

• We choose a generic direction CERd and orient the edge-graph accordingly



e define 
$$h_i^c := \# \text{ vertices with out-degree} = i$$

$$\rightarrow$$
 we show that  $out(v) := \# edlser pointing out of volumes these numbers are independent of c!$ 

· we express the f-vector in terms of the hi.

$$f_{k} = \sum_{\substack{v \in \mathcal{F}_{0}(P) \\ \text{out}(v) \geq k}} \binom{out(v)}{k} = \sum_{i=k}^{d} \binom{i}{k} h_{i}^{c}$$

- 1. every face of P contains a unique minimal vertex wrt. The orientation

  7. every set of k edges emanating from a vertex belongs to a unique face (only true for simple polytoper)
  - · we show that therefore the his are also determined by the f-vector, independent of c.

• define 
$$\overline{f}(t) := \sum_{i=0}^{d} f_i t^i$$
 and  $H(t) := \sum_{i=0}^{d} h_i^c t^i$ 

$$\rightarrow \mathcal{H}^{c}(t+1) = \sum_{i=0}^{d} h_{i}(t+1)^{i} = \sum_{i=0}^{d} h_{i} \sum_{k=0}^{i} {i \choose k} t^{k}$$

$$= \sum_{k=0}^{d} t^{k} \sum_{i=k}^{d} {i \choose k} h_{i}$$

$$\stackrel{(x)}{=} \sum_{k=0}^{d} t^{k} f_{k} = \mathcal{F}(t)$$

$$\rightarrow$$
 H° and the h° independent of c and we can write H and h; instead.

$$\overrightarrow{Def}: h - vector \overrightarrow{h} := (h_0, ..., h_{cl})$$

- · super important in modern polytope theory
- · the "right way" to look at fore numbers
- o appears everywhere: reconstruction from the edge-graph -computing volumer (thrhavt theory) -shellings ...
- · Since the h; are independent of the direction vector

$$h_i = h_i^c = h_{i-i}^c = h_{d-i}^c = h_{d-i}$$

- 1. if i edger are going out of v then d-i edges are coming in.
- 7. C-+-c flips in- and out-going eager

$$\rightarrow$$
  $h_i = h_{d-i}$  Dehn-Sommervike equations in h-basis.

• note that H(t) = F(t-1), therefore

$$h_k = \sum_{j=k}^{cl} (-1)^j {i \choose k} f_i$$

which yields the equations in f-basis.

#### Remorks:

- \* k=0 gives the Euler-Poincaré identity for simple / simplicial polytopes
- the hi count something, therefore hi≥0 and we get potentially useful inoqualities:

$$h_0 = 1 \ge 0$$
 trivially true  
 $h_1 = fal_{-1} - al \ge 0$  easily verified  
 $h_2 = fal_{-2} - (al_{-1})fal_{-1} + (al_{-2}) \ge 0$  ... not so trivial  
anymore

· much stronger inequalities are known

Thm: the h-vector is unimodal
i.e. its entries increase, and then decrease
with only one peak.

This is a deep theorem (see g-theorem later)

- If a polytope is  $> \lfloor d/2 \rfloor$  -neighborly then more then half of the entries of f-vector as see with simplex
  - -> f-vector is the same as simplex
  - $\rightarrow$  Ex: it is a simplex
- We can compute (ull f-vector of cyclic polytope  $C_{\alpha}(n)$   $-k \leq \lfloor \frac{d}{2} \rfloor : \quad f_{k} = \binom{n}{k+1} \qquad \qquad \text{(not included in the lecture)}$

$$\int_{k} = \frac{n - \delta(n - k - 2)}{n - k - 1} \sum_{j=0}^{\lfloor a/2 \rfloor} {n - 1 - j \choose k + 1 - j} {n - k - 1 \choose 2j - k - 1 + \delta}$$

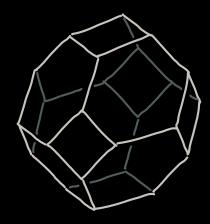
where  $\delta = 0$  if d is even, and  $\delta = 1$  otherwise.

$$\rightarrow f_{d-1}(C_{\alpha}(n)) \in O(n^{d/2})$$

The Denn-Sommorville equations are the only linear relations between the face numbers of simple / simplicial polytopes

Applying the Dehn-Sommorville equations to particular polytopes for which is known that the face-numbers encode important combinatorial sequences can yield noval relations! (not included in the lecture)

#### Example



permutanedron (n!)



(Catalen numbers)

## 4.3. The upper bound theorem & g-theorem

- e given a d-polytope with n vertices, how many k-faces can it have?
- o not only is this upper bound known, there even exists a polytope which attains this bound for all k!

-> the cyclic polytope Cd(n)

Thm: (upper bound theorem)

If P is a d-polytope with n vertices then

$$f_k(P) \leq f_k(C_d(n))$$
 for all k

- This is one of the big theorem s. Combinetorial proofs are known, but the modern approaches are via commutative algebra.
- of -vectors of simple/simplicial polytopes are completely charactenized

$$f$$
-vector  $\mapsto$  h-vector  $\mapsto$   $g$ -vector  $g_0 := 1$ ,  $g_k := h_k - h_{k-1}$ 

(Actually: the g-vectors got clossifical)

Thm: An f-vector belongs to a simplicial polytope iff its g-vector is an M-sequence.

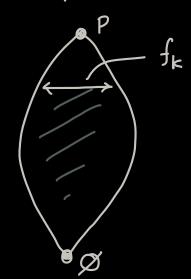
M-sequence means 
$$g_o = 1$$
 and  $g_{k-1} \ge \partial^k(g_k) \ge 0$ 

where  $\partial^k: \mathbb{N} \longrightarrow \mathbb{N}$  is some not too complicated function.

only provon 2018 by Karim Adiprosito using some quite advanced algebraic techniques.

# 4.4. Other facts about face numbers

• is this picture of the face lattice accurate?



Q: is the f-vector unimodal?

NO and for general

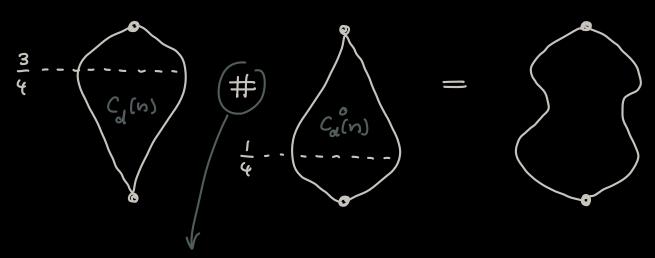
polytoper counterexamples

can be constructed

"quite earily."

#### Example:

o f-vector of cyclic polytope  $C_d(n)$ peaks at  $n^3/4d \rightarrow dual peaks at <math>n^4d$ 



connected sum (gluing polytopes along faces) adds up the f-vectors (elmost).

yields counterexamples with d=8 and  $f_0 \sim 7000$  or d=9 and  $f_0 \sim 1500$ 

- · Ca(n) # Ca(n) is neither simple nox simplicial
- in fact, it is known that for  $d \leq 15$ simple I simplicial polytoper have a unimodal f-vector.
- BUT: Using the g-theorem counterexamples can be found with d=20,  $f_0=169$  or d=30,  $f_0=47$
- on top or bottom?

$$f_k \stackrel{?}{=} min (f_0, f_{\alpha-1})$$
 (Báránys conjecture)

-> just solved this year (2022)! by Joshua Hinmen

• Shape of 4-dimensional f-vector can be measured more precisely

fatness (P) := 
$$\frac{\int_1 + \int_2}{\int_0 + \int_3}$$
 = thickness "in the middle" bottom"

OPEN: Is the fatness of 4-polytoper bounded?

· Highest known fatness is 5+E

OPFN: Does there exist a 4-polytope all whose 3-faces are icosaheava?

 $\rightarrow$  if yes, this polytope would have an outstandingly large fot near (=0.2)

- A lot of ongoing research tries to generalize neighborhyness, Dehn-Sommerville, upper bounds and characterizations to other classes:
  - centrally symmetric
  - cubical (all faces are cubes)